

## Search

### Agents

Agents interest of  
complete process  
(simple system)

Agents interest of what  
will happen after ending  
the process of find  
best solution.  
(complex system)

#### ① Reflex agent (simple system)

- a) choose action based on current percept.
- b) may have memory or model of world's current state.
- c) Do not consider future consequence of their actions.  
"Do not ask ~~if~~ what if"
- d) Consider how the world is (Cannot care about future actions)
- e) Can it be rational (yes)?  
↳ yes, if it is provided by if conditions.

## 2] Planning Agents (complex system)

- a) Ask "what-if"
- b) Decisions based on (hypothesized) consequences of actions.
- c) must have model of how world evolves in response to actions.
- d) must formulate a goal (test)  $\Rightarrow$  must reach the goal.
- e) consider how the world would be.

Planning	Replanning
→ Agent make Plan according to search.	→ Agents make Plan for the first step according to search then take decision to complete first step. <u>then</u>
→ It take period of time to plan then take decision according to plan.	→ It replan for next step according to new search and take new best decision to start work for next step.
→ Plan only one time	→ Plan for each step ↳ no. of plans = no. of steps.

## \*optimal vs complete Planning

↳ Agent, not only reach goal or take action but also it want to get optimal solution of problem. (Search for best way to goal or minimum time to reach goal)

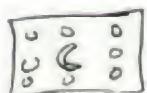
## \* Search Problems:

↳ process to take right decision.

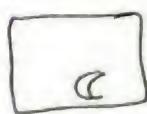
It consists of

1) state space (world): All possible cases of the model.

ex mouse want to eat 9 piece of cheese.



→ mouse eat piece on center



→ eat the last piece.

2) A Successor Function (actions, costs)

↳ For any state what action can I take and what cost of taking it.

### 3) A start state an goal test

start state: state from which i can start working.

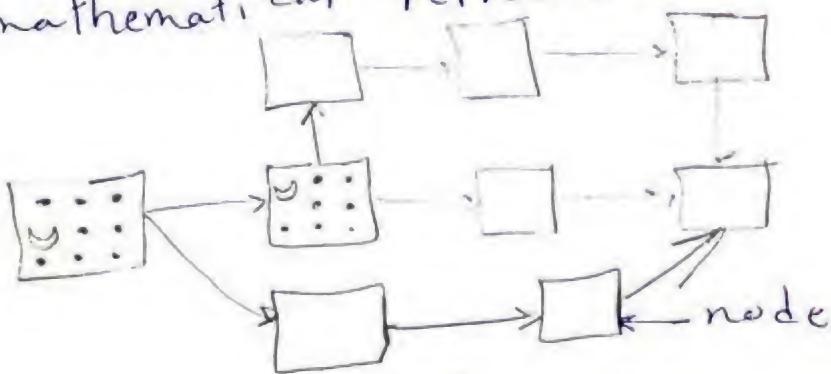
goal test : the final state.

→ solution of search problem is sequence of action which transforms start state to goal state.

### \*Uniformed search methods

#### ① State space representation

→ mathematical representation of search Problem.



• Nodes → world configurations.

• Arc → successors (action results)

• goal test is set of goal nodes.

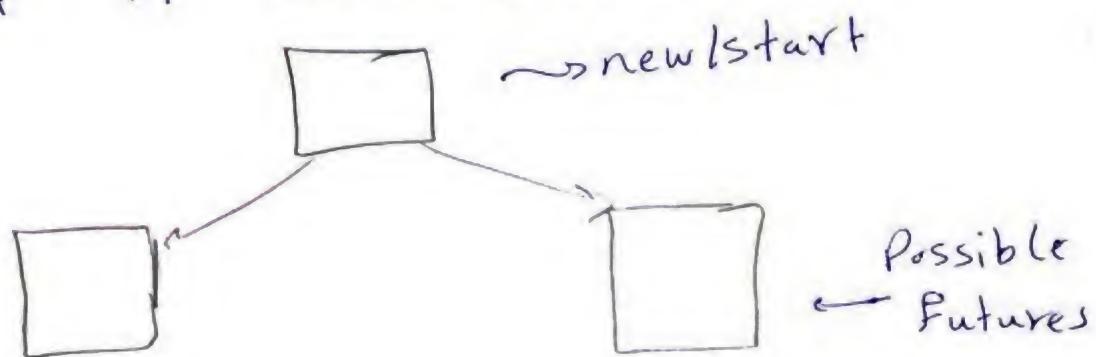
→ here we represent needed states only.

→ we can rarely build this full graph.

→ there is no start state. but goal test.

## ② Search trees

↳ part of search graph.



### Search tree

→ start state is the root node.

→ children correspond to successors.

→ we can never actually build whole tree.

### \* searching with search tree

- try to expand as few tree nodes as possible
- maintain a fringe of partial plans under consideration.
- try to expand as few tree nodes as possible

(General Tree  
Search)

Fringe : all of plans that may yet work.

→ all ways that can reach me from start to goal during world state search tree.

→ one of this frings is my solution (optimal solution)

### Expansion

↳ picking some thing out of the fringe and if it is not a goal already.

→ process of expanding new state from current state on one of frings. we can't do this process if we reach goal.

### \* Exploration strategy

↳ what fringe nodes do you explore next?  
↳ what will happen next?

### \* Main Question:-

↳ which fringe nodes to explore to reach goal?

### Depth-First Search

strategy: expand deepest node first

Implementation: fringe is LIFO stack.

Solution: left most solution.

## Search Algorithm Properties

→ For any Algorithm we check For 4 Properties:-

1] Complete:

↳ Ability to reach goal if it is exist.

2] Optimal:

↳ Ability to reach to best solution.

3] Time Complexity

↳ time to reach to solution

↳ time to expand all nodes of one fringe equal to time to expand one node of fringe \* no. of fringe nodes.

4] Space Complexity

↳ size of fringe in memory - size of nodes of fringe in stack.

\* no. of nodes in entire tree.

$$1 + b + b^2 + \dots + b^m = O(b^m)$$

~ branching factor:- start node can make a branch for b number of children.

## Depth First Properties

### 1] Time complexity

- what nodes DFS expand?
  - ↳ some left prefix of tree.
  - ↳ could process the whole tree!
- ↳ worst case expand all nodes (right most)
  - ↳ best case expand no nodes (start  $\equiv$  goal)

worst case:  $O(b^m)$ , best case =  $O(1)$

### 2] Space complexity

- ↳ only has siblings on path to root, so  $O(b^m)$
- ↳ worst case we get solution on last tier.

### 3] Complete

- ↳ m could be finite, if we prevent cycles.
- solution can only be found if
  - 1. it exists
  - 2. Finite Algorithm.

### 4] Optimal

- no, it finds left most solution regardless of depth or cost.

## 2] Breadth -First (BFS)

strategy: expand a shallowest node First.

Implementation: fringe is FIFO queue.

solution: shallowest solution (shortest fringe)

\*

### \*DFS & BFS

→ BFS will outperform DFS when:

- 1) need less time complexity.
- 2) need shallowest solution.
- 3) need optimal solutions for costs equal to 1.
- 4) need complete solutions.

→ DFS will outperform BFS when:

- 1) most solutions at left side of tree.
- 2) all ~~solutions~~ at the last level.
- 3) need less space complexity.



### 3) uniform cost

strategy: expand a cheapest node first.

Implementation: fringe is priority queue

solution: cheapest solution.

→ expand node of cheapest code.

#### \*Advantages

→ complete and optimal.

#### \*Disadvantages

1. no information about goal location.

(expand) after knowing (goal)      Nodes, not node ←  
· visit with node in { }

2. explores options in every direction.

↳ no determined direction on our work.

## \* The one queue : Priority queues

- conceptually, all frings are priority queues.
- For DFS and BFS you can avoid the  $\log(n)$  overhead from actual priority queue with stacks and queues.
- Can even code one implementation that takes variable queuing object.

## Informed search

↳ we need to know information about goal location to know if I work on correct way or not.

→ we have one function, 2 search Algorithms:-

### 1) Heuristics

↳ function takes state of state space, and give number which represent how far the goal location from my location to doing process.

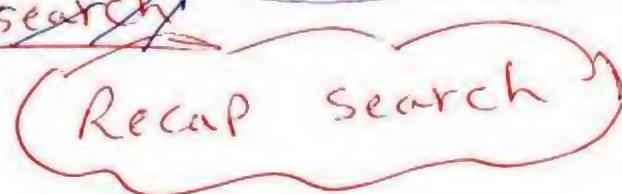
## \* Greedy search

- ↳ Search Algorithm use idea of Heuristic.
- ↳ It is not an optimal search Algorithm.

## \* A\* Algorithm

- ↳ It collects all last search Algorithms ideas to get very good search Algorithm.

## \* Graph search



### 1) Search Problem

- a) states (configuration of world)
- b) actions & costs.
- c) successor function (say how states respond to actions)
- d) start state & goal state.

### 2) Search tree

→ nodes: Plans for reaching states.

- a) nodes: Plans for reaching states.
- b) plans have costs (sum of action costs)

### 3) Search Algorithm

→ systematically builds search tree.

- a) systematically builds search tree.
- b) chooses an ordering of the fringe.

- c) optimal: find the least-cost plans.

## Ex: Pancake Problem

Problem: need to arrange pancake from big to small.

states: shape of pancakes during flipping to reach goal.

costs: no. of pancakes flipped.

start state: 

goal state: 

Algorithm: we can use UCS algorithm according to cost also we can use DFS or BFS.

## Search Heuristics

heuristic is

- a) Function that estimates how close a state to goal.
- b) Designed for particular search problem.
- c) we make it every step to see if ~~some~~ is close to goal or not.
- d) every search problem need different heuristic

According to nature of problem.

## Simple Heuristics

Heuristic: no. of largest pancake that is still out of place.

## Greedy search

strategy: expand the nodes that seems closest to goal according to heuristic.

Heuristic: estimate the distance to nearest goal for each state.

Implementation: Priority queue :-  
solution: best heuristic solution.

→ It doesn't care about cost.  
→ cares only about heuristic

~~common case~~: Best-First takes you to wrong goal.  
worst case: like badly-guided DFS.

solution → combine UCS and Greedy search.

## A\* search

→ combining UCS and Greedy Search.

\* uniform cost orders by path cost or backward cost  $g(n)$ .

\* Greedy search is by goal proximity or forward cost  $h(n)$ .

↳ should we stop when we enqueue goal?

↳ No only stop when we dequeue goal.

### Admissibility

- \* Inadmissible (Pessimistic) heuristics break optimality by trapping good plans on fringe.
- \* Admissible (optimistic) heuristics slow down bad plans but plans never outweigh true costs.

### (Admissible heuristic)

→ A heuristic  $h$  is admissible if:

$$0 \leq h(n) \leq h^*(n) \rightarrow h^*(n) = g(n)$$

→  $h^*(n)$  is true cost to nearest goal

$h(n) > h^*(n) \rightarrow$  (<sup>optimal</sup> solution) no update

### \* A\* applications

\* video games. \* language analysis.

\* Machine translation \* speech recognition.

\* Robot motion Planning.

\* resource planning problems.

## \* Creating Admissible heuristics

most of work in solving hard search problems optimally is in coming up with admissible heuristics.

→ often, admissible heuristics are solutions to relaxed problems where new actions are available.

→ inadmissible heuristics are often useful too.

### Ex 8x Puzzle

7	2	4
5		6
8	3	1

start state

3	7	1
2	4	5
E	R	6

Action

	1	2
3	4	5
6	7	8

Goal state

states? all cases that i can move any number from its location to another one if it is next to free space

How many states →  $8! =$

actions? moving number to free space next to it.

- \* no. of successors from start state?
- maximum 4 moves (minimum 2 moves)
- \* Cost should be? 1 every time
- \* Heuristic: no. of tiles misplaced.

$$h(\text{start}) = 8$$

- \* Why it is admissible?

1) There is relaxed-problem heuristic  
(easiest solution but need more work)

2) Direct move  
↳ move every number from start state to final state.

$h = \text{no. of moves}$ .  
 $h(1) = 3$  (moves to reach to its correct state at goal state)

$$h(2) = 1, h(3) = 2 \text{ and so on.}$$

∴ Total  $h = 18$  from start to goal.

Admissible if  $0 \leq h(n) \leq 2(n)$

→ How about using actual cost as heuristic?

$$f(n) = h(n) + g(n) = 2g(n) = 2h(n)$$

$$\text{if } g(n) = h(n)$$

↳ we go back to UCS by double cost

value trade off: complex work, neglect heuristic

with A\* → trade off - between quality of estimate and work per-node.

### \* Trivial heuristic

• bottom of lattice is zero heuristic

• top " " is the exact " if  $(h=0)$  no heuristic  $\Rightarrow$  go back to  $\begin{cases} \text{uniform} \\ \text{cost} \\ \text{search} \end{cases}$

Dominance:  $h_a \leq h_c$

$\forall n: h_a(n) \geq h_c(n)$  for all nodes

we choose  $h_a(n)$  (best one)

⇒ Heuristics from a semi-lattice

↳ max of admissible heuristics is admissible

$$h(n) = \max(h_a(n), h_b(n))$$

we choose

## \* Graph search

↳ failure to detect repeated states can cause exponential more work.

idea: never expand state twice.

### Implementation:

- \* tree search + set of expanded states (<sup>closest</sup> set)
- \* expand search tree node by node, but never expand node twice.
- \* Before expanding a node, check if it is expanded before or not.
- \* if not now, skip it, if new add to closest set.

### Important

- \* store closest set as a set, not a list.
- \* complete ~ if goal exists, we will find it.
- \* optimal ~ no.

(19)

	Depth First search	Breadth First search	Uniform Cost Search
nodes expanded	<ul style="list-style-type: none"> <li>→ some left prefix of the tree.</li> <li>→ Could process the whole tree.</li> <li>→ if <math>m</math> is infinite, takes time <math>O(b^m)</math></li> </ul>	<ul style="list-style-type: none"> <li>→ processes all nodes above shallowst solution.</li> <li>→ let depth of shallowst solution be <math>s</math>.</li> <li>→ search takes time <math>O(b^s)</math></li> </ul>	<ul style="list-style-type: none"> <li>→ process all nodes with cost less than cheapest solution.</li> <li>→ takes time <math>O(C^*/\epsilon)</math></li> <li><math>C^*</math> → sol. costs</li> <li><math>\epsilon</math>.</li> </ul>
space that strings take	$O(b^m)$	$O(b^s)$	$O(b^{C^*/\epsilon})$
complete	only if we prevent cycles	$s \rightarrow$ finite Yes	Assume best sol. has finite cost and min. $\epsilon$ $\Rightarrow$ Yes
optimal	No	only if costs are all 1	Yes

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→ heuristic

\* Function estimates how close a state is to a goal.

\* Designed for particular search problem.

\* Ex: Manhattan distance.

### Greedy search

\* strategy expand a node that you think is closest to a goal state.

↳ heuristic → estimate of distance to nearest goal for each state.

\* Common case

↳ best-first takes you to wrong goal.

### A\* search

↳ only stop when we dequeue a goal.

#### Admissible Heuristics

↳ heuristic  $h$  is admissible (optimistic) if

$$0 \leq h(n) \leq h^*(n)$$

where  $\Rightarrow h^*(n)$  → is true cost to nearest goal.

## UCS vs A\* Contours

UCS → expands equally in all "directions".

A\* → expands mainly toward the goal but does  
hedge its bets to ensure optimality

## A\* Applications

video games, language analysis, speech ~~recognition~~,  
recognition, machine translation, Pathing.

→ search problem consists of

1) state space

2) successor function  
(with actions, costs)

3) start state and goal test.

solution → is sequence of actions (a plan)  
which transform start state to goal state.